

Is QoS Necessary?

QoS Mechanisms vs. Bandwidth Provisioning

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Introduction

- Overview of QoS Mechanisms
- Evaluation of mechanisms vs. bandwidth
- Deployment experience
- Conclusions

Defining QoS: Application perspective

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 QoS is defined in terms of loss rate, throughput, delay

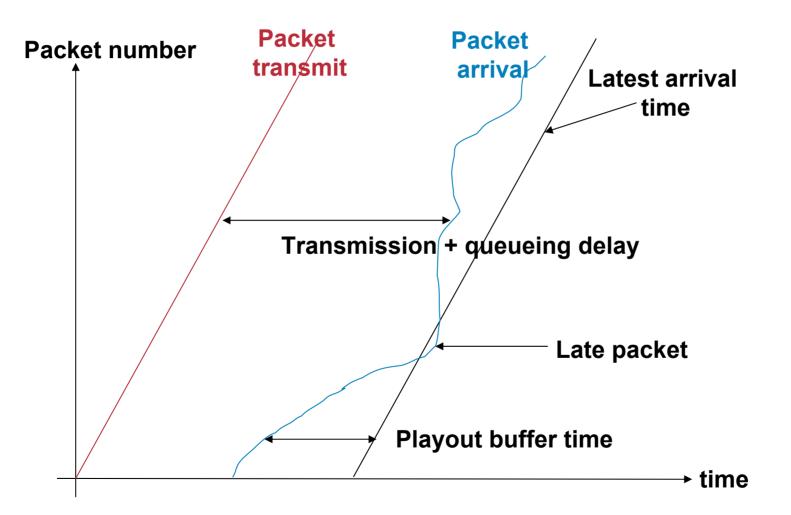
Throughput & loss rate tightly coupled for many apps (e.g. TCP-based)

Applications that need timely delivery of packets (e.g. Voice) sensitive to *delay distribution*

e.g. 99th percentile delay

"Late" packets typically useless – "early" packets can be buffered

QoS for multimedia applications

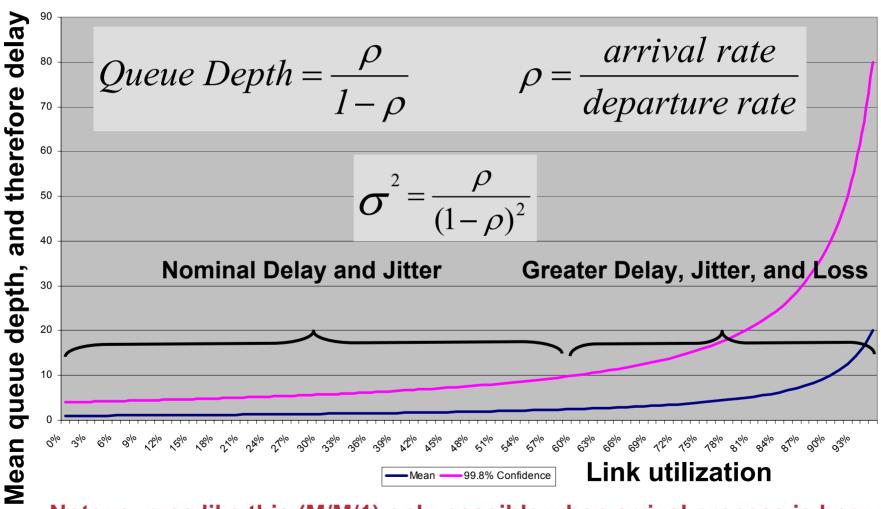


Defining QoS: Mechanism perspective

- If bandwidth is plentiful, packets are never delayed or lost, so QoS needs of apps are met
- QoS often defined in terms of mechanisms deployed when bandwidth isn't plentiful

Mechanisms include classification, policing, shaping, marking, queueing, differential dropping

Queuing Theory and QoS



Note: curves like this (M/M/1) only possible when arrival process is known

Throwing bandwidth at the problem

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- As link bandwidth $\rightarrow \infty$, utilization $\rightarrow 0$, so delay, loss $\rightarrow 0$
- Problems:

As bw $\rightarrow \infty$, cost also $\rightarrow \infty$

When does the cost/quality tradeoff work?

Pent-up demand frequently drives utilization up as bw becomes available

e.g. TCP seeking to congest the bottleneck link

e.g. video becomes more attractive as bw increases

"Arms race" to see who becomes the bottleneck

The growth of bandwidth

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 Fashionable to talk about "effectively infinite" bandwidth solving all problems

WDM, Photonic networks, OC-n SONET/SDH

Used by research, Internet backbones, and large corporations

Corporate practical reality

T-1 remains the most common access link in the US

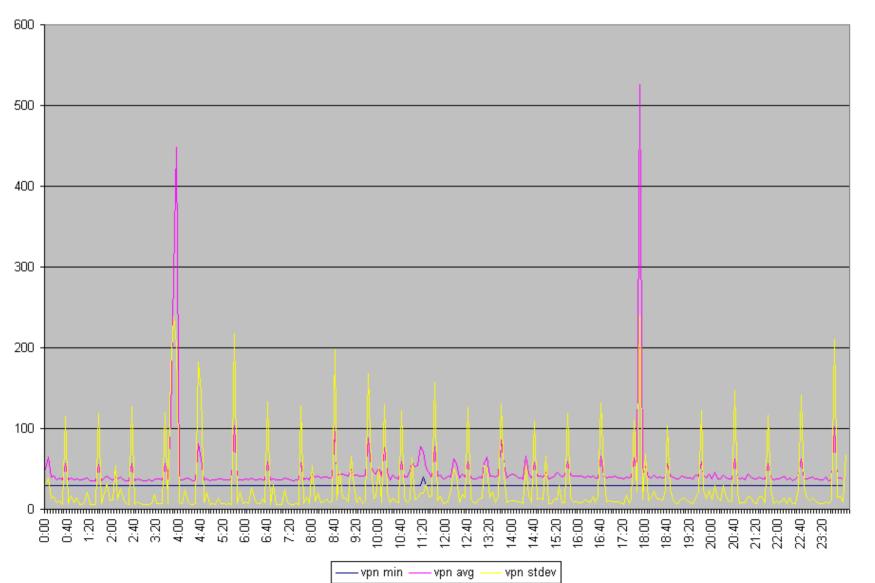
Globally, 64-128 KBPS remains a "large" access bandwidth, and more is costly

Consumer practical reality

"Broadband Revolution" fitful and difficult to make money on Perhaps half of Internet access still via dialup Even broadband connections prone to congestion

Broadband access delays

Data courtesy of Fred Baker



TCP behavior

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- 95% of Internet traffic today runs over TCP
- The fact that the Internet hasn't collapsed can be traced to TCP congestion avoidance & "slowstart"

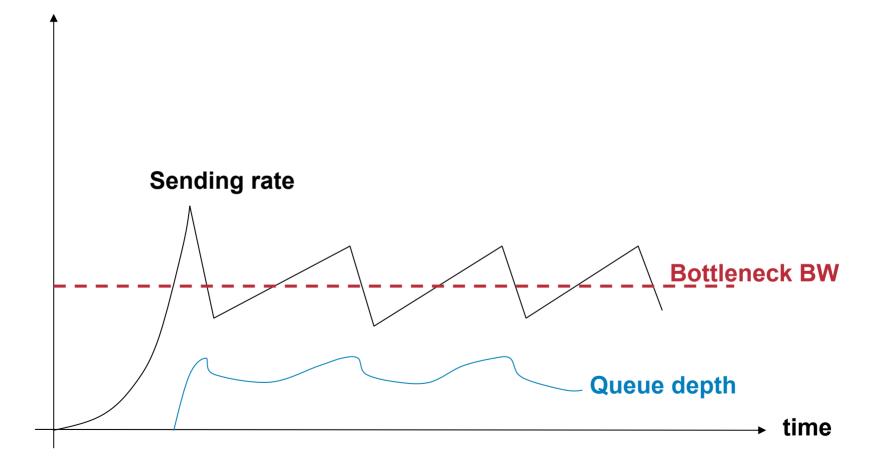
Q. What happens if TCP becomes less dominant?

TCP "avoids" congestion by causing it, then reacting

Attempts to find BW of the "bottleneck link" and send at that rate

TCP behavior

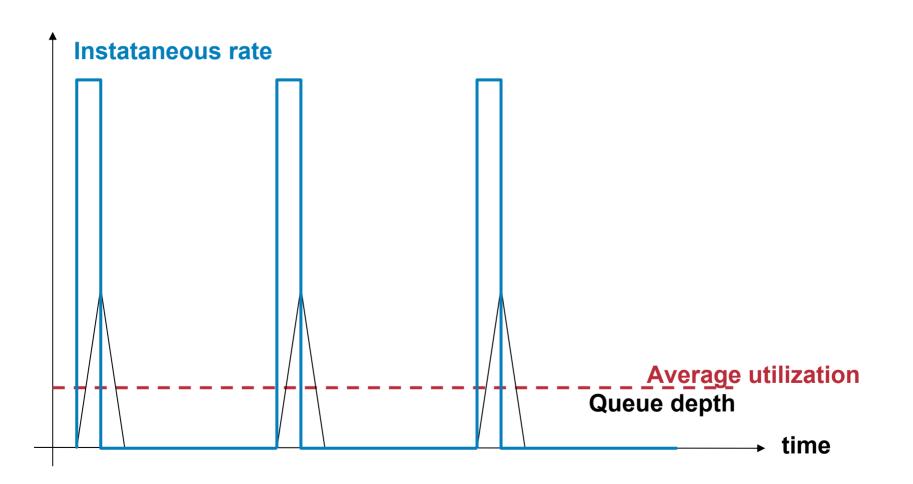






- A 10ms delay per hop may be too much for voice
- It takes < 10ms to create a 10ms queue ⇒Traffic bursts at the millisecond level can matter
- Average utilization measured over seconds or longer doesn't capture this

Timescales



QoS mechanisms

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 If BW allocation is the only tool, then need to provide enough BW to meet the most stringent QoS needs of any app

provide low delay to all traffic, not just voice/video

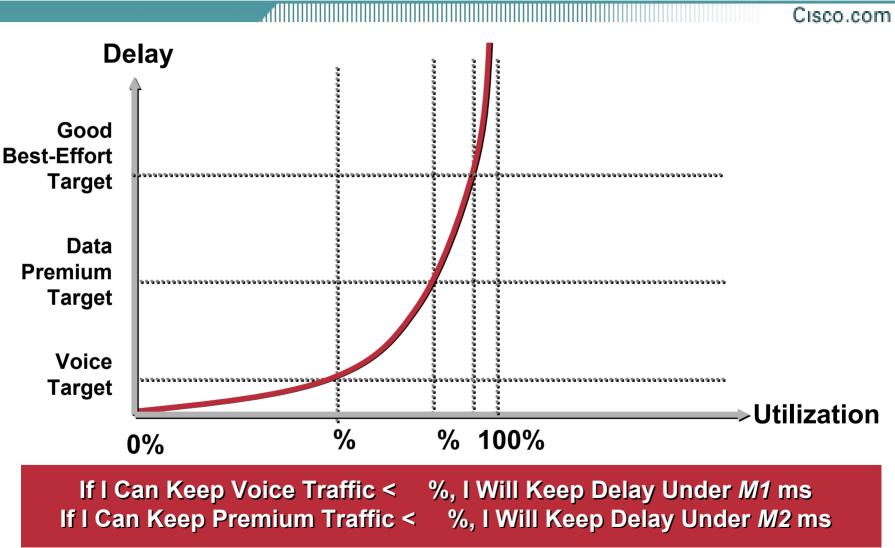
low loss to all apps (maybe overkill for voice)

 Most QoS mechanisms aim to allocate the delay & loss to apps that are less sensitive to it

Can't create BW, just control who gets access to it and when

Sorting traffic into distinct queues creates different utilizations for each class

Customizing Delay/Load Trade-Off



QoS mechanisms – other reasons

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Service differentiation

Providers would like to extract money from customers willing to pay, e.g., corporate users with mission critical data, voice

Requires mechanisms to make some traffic see "better" service

Protection among classes

Interaction between bursty & non-bursty sources Interaction between TCP-friendly & others Cisco.com

QoS Mechanisms

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QoS Mechanisms

- Classification: sorting packets into categories
- Marking
- Queueing/scheduling: determining which order packets get sent
- Policing/metering: counting packets & taking some action (e.g. remarking, dropping)
- Shaping: limiting the bandwidth to a class
- Dropping
- Admission control

Classification

- Sort packets into categories

 e.g. voice, video, premium, best effort
 Can use almost any criteria
 packet contents IP header, application data
 point of entry (representing a customer)
 Per-flow or aggregate classification
- Necessary precursor to most other QoS functions

Queueing/scheduling

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- Determine which order packets get sent
- Determines how much bandwidth a class receives
- Basic approaches:

Priority – send higher priority traffic always

Creates the illusion that other traffic is absent, i.e. emulates a less loaded network

(Weighted) Fair Queue – like round robin, allocates a share of the link to a class

Enables utilization to be controlled on per-class basis

Work-conserving algorithms avoid bandwidth wastage

Policing/Metering

Count the rate of packet arrivals in a class

Typically uses a "token bucket" with a rate and burst size

Tokens accumulate at some rate, up to max burst, and are "spent" as each packet arrives

If packet arrives at empty bucket, either drop it or downgrade its marking

Provides a means to limit the number of packets in a class

Limits can be on a per-customer basis (or finer)



- Traditionally, packet drops occurred only when queues were exhausted (tail-drop)
- Active queue management drops more intelligently
 - Send the message to a TCP app with just one packet drop, not several
 - Drop preferentially from less sensitive/less important applications
 - Drop preferentially from traffic that was previously policed as "out of contract"

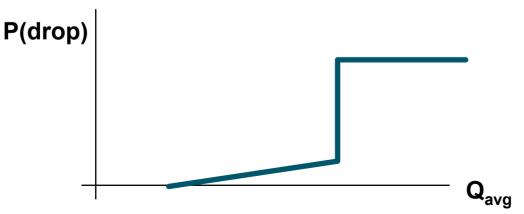
Random Early Detection

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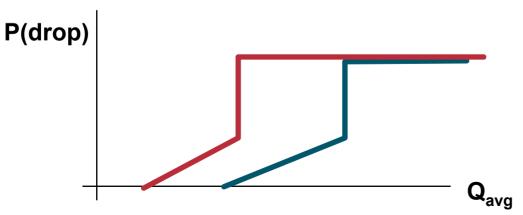
 Average queue length is monitored to detect onset of congestion

Averaging time constant is O(RTT)

 As small percentage of packets is dropped, TCP backs off, congestion is averted



- Just like RED, but with different drop profiles for different classes
- Can carry extra traffic opportunistically & shed load as needed



Explicit Congestion Notification

- A simple enhancement to RED
- Packets are marked rather than dropped
- TCP congestion avoidance responds as if drop occurred other transports may also react appropriately
- Congestion avoidance without loss
- Host participation required

Packets marked "ECN capable"

Receiver conveys marking back to sender at transport layer

Admission Control

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- All mechanisms so far are in "forwarding plane" – admission control is the "control plane"
- Determining ahead of time if the resources are available

Avoiding the "everyone is high priority" syndrome

Most valuable if the application would prefer no service to impaired service, e.g. "fast-busy" for voice

Better to have 99 good quality calls and 1 fast busy than 100 poor calls

Applications with steep "utility functions" the best candidates

Focusing on IP QoS

 IP the dominant networking protocol today Global reach

Heterogeneity of link layers and applications

- IP QoS is not an oxymoron (as some had thought)
- Mechanisms are in fact generic

Differentiated Services Overview

- "Diff-serv" is a set of Internet standards
- Clearly the preferred QoS technology for ISPs & enterprises today
- Near-minimal complexity

e.g. can deploy DS with just 1 header bit and 2 "per-hop behaviors" (PHBs)

 Edge behavior (classification, marking, policing etc.) + core behavior (PHBs) provides services variety of services from a single PHB

Diff-serv Contributions

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 Standardized definition of the "Diff-serv Code Point" (DSCP) in IP header

(after years of confusion about TOS)

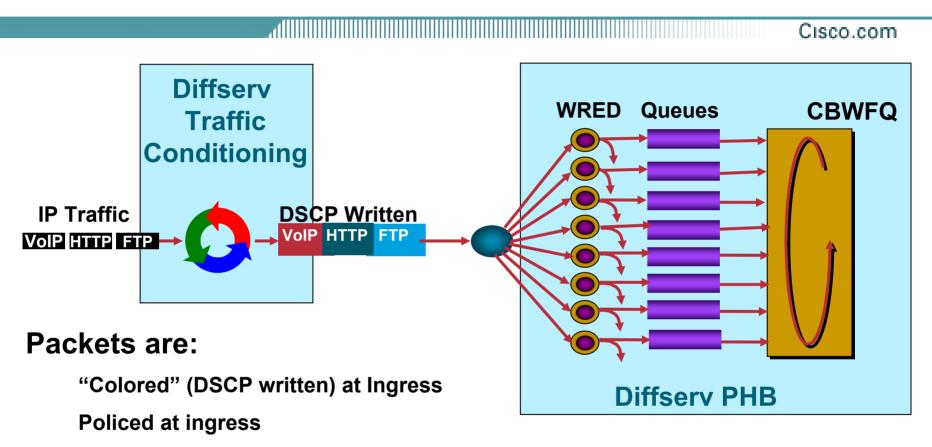
Small set of standard PHBs

Expedited Forwarding (like a priority queue)

Assured forwarding (WFQ & WRED)

 An overall architecture for DS Mostly formalizing ideas already in use

Putting it all together



Classified before queuing

Potentially discarded by WRED

Placed in queues based on DSCP

Scheduled by CBWFQ (and/or LLQ)

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 Classifier + token bucket policer at network edge

recognize & meter traffic in need of isolation; set DSCP = x

- Dedicate a queue (and some bandwidth) to DSCP = x
- Effect is to run this traffic on its own logical network (with controlled utilization)

Example Service (2)

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 Customers buy an SLA involving an allowed amount of "premium" data service

Specified as a token bucket

- Policer at provider ingress marks traffic below the rate as "premium", excess is marked as "best effort"
- These markings are used by WRED to ensure that excess traffic is sent if possible, but dropped under congestion
- Note that WRED is implemented on a single queue, so no intra-flow misordering

Good news for TCP



- It was widely perceived that IPv6 would enhance the QoS of IP
- In fact, QoS mechanisms for IPv6 are identical to those of IPv4

Both use 6 bits in header for Diffserv, 2 bits for ECN

Both supported by RSVP

Evaluating Mechanisms vs. Bandwidth

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Evaluation – It's all about cost

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Bandwidth comes at some price

Very much dependent on environment (LAN/WAN, last mile/backbone, own/lease, regulatory factors)

Or not...for many consumers, higher speed access isn't there at any price

QoS mechanisms incur operational costs

Train personnel to configure correctly

Negotiate and bill for (more complex) SLAs with customers

Monitor network for incorrect or outdated configurations, risk of SLA violations

Possible performance impact on equipment

Other Cost/Benefit Factors

 If QoS mechanisms make more applications work on common infrastructure (e.g. VOIP) then there are cost savings

The driving reason behind enterprise VOIP is avoiding toll costs, leveraging existing capacity

Benefits of traffic isolation

 Separating "premium" customers from others makes stronger (and more lucrative) SLAs possible

e.g. may continue to deliver SLA even during failure/overload scenarios

Separating traffic types can be good for all

TCP performs worse when sharing queues with "nonresponsive" flows

Voice performs worse when sharing queues with more bursty traffic types

Asian perspective

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Application needs

Probably not too different from anywhere else

VOIP more attractive wherever TDM voice is costly relative to IP bandwidth

Link costs

Glut of capacity less apparent than in the US

Transoceanic links inherently costly (and timeconsuming to deploy/upgrade)

Regulation frequently drives costs up

Providers who don't own their links much more sensitive to cost

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Deployment Experience

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QoS Deployment

- Backbone ISPs
- VPN Service Providers
- Enterprise

- Most today have built in sufficient capacity to make QoS mechanisms superfluous
- Hard to start selling premium services when Your basic service is already good enough for voice So is your competitors'
- Some are carrying VOIP and using basic priority queueing mechanisms
- Interest in QoS mechanisms higher when failures are considered
- Interest may rise as capacity glut subsides

VPN service providers

- Over 100 SPs today offer "Layer 3" VPN services (MPLS/BGP VPNs)
- Mostly selling as a frame-relay replacement
- High customer demand for "CIR"-like services
- Many SPs using PQ for VOIP and WRED for premium data services



- Bandwidth in the campus is cheap enough to overprovision
- Inter-site links are the bottlenecks
- Plenty of interest in PQ mechanisms for VOIP Admission control often needed also
- Mission critical apps with response-time requirements also common
- Video getting more common

Conclusions

In theory, bandwidth could solve all QoS problems

 In practice, plenty of need for QoS mechanisms BW too expensive/not available

Protection from "greedy" applications

Protection against failures

Premium services to enhance revenue

 There are many useful IP QoS technologies deployed today